

Slide 3: Recap on Logical Implication (Entailment) \models

- Entailment notation: $p \models q$ if and only if the implication $p \Rightarrow q$ is a tautology.
- Example:
 - $p \wedge q \models q$
 - Truth table for $p \Rightarrow q$:

p	q	$p \wedge q$	$p \Rightarrow q$
T	T	T	T
T	F	F	F
F	T	F	T
F	F	F	T

Slide 4: $(r \Rightarrow s) \wedge (r \Rightarrow \neg s) \models \text{false}$

- Intuitively, if r implies both s and $\neg s$, then r must be false.
- Truth table for $(r \Rightarrow s) \wedge (r \Rightarrow \neg s)$:

r	s	$\neg s$	$(r \Rightarrow s) \wedge (r \Rightarrow \neg s)$
T	T	F	F
T	F	T	F
F	T	T	T
F	F	T	T

Slide 5: $p \vdash q$

- Notation: $p \vdash q$ means q is provable from p using inference rules.
- Example:
 - $A \Rightarrow B, \neg A$, therefore $\neg B$

Slide 6: Differences Between \models and \vdash

- \models indicates semantic entailment (truth conditions).
- \vdash represents syntactic derivation (inference rules).

Slide 7: Recap on Inference Rules

- Example inference rules:
 - Modus Ponens (\Rightarrow Elim):
 $p \Rightarrow q, p \vdash q$
 - Conjunction Introduction (\wedge Intro):
 $p \vdash q, p \vdash r \vdash p \wedge r$
 - Conditional Proof (\Rightarrow Intro):
 $p \vdash r, p \vdash s \vdash p \Rightarrow (r \wedge s)$

Slide 8: Layout of an Inference Rule

- Premises above the line, conclusion below the line.
- Example inference rule (\Rightarrow Intro):

$$p \vdash r, p \vdash s \quad p \Rightarrow (r \wedge s)$$

Slide 9: Presentation of Proofs

- Steps:
 - Number each step.
 - Justify each step with previous line(s) and inference rule used.

Slide 10: Deriving $\neg p \Rightarrow r$ From $(p \wedge q) \vee r$

- Example proof:
- $$(p \wedge q) \vee r, \neg E \dots \neg p \Rightarrow r$$

Slide 11: Two Special Inference Rules

- Deductive Theorem (\Rightarrow Intro):
- $$p \vdash r, p \vdash s \quad p \Rightarrow (r \wedge s)$$
- Reductio ad absurdum (\neg Intro):
- $$p \vdash r, p \vdash \neg s \quad p \vdash \neg r$$

Slide 12: Conditional Proofs

- Strategy: Assume p, deduce q if possible, discharge assumption.
 - Example:
- $$(p \wedge q) \vee r \dots \neg p \Rightarrow r$$

Slide 13: Indirect Proofs

- Strategy: Assume negation of goal, deduce contradiction.
 - Example:
- $$(p \wedge q) \vee r \dots \neg p \Rightarrow r$$

Slide 14: Solution to Exercise

Given argument: A (You eat carefully) B (You have a healthy digestive system)
 C (You exercise regularly) D (You are very fit) B D E (You live to a ripe old age) $\neg E$ Therefore, $\neg A \quad \neg C$

Proof:

Line	Formula	Justification
1	A B	Premise
2	C D	Premise
3	B D E	Premise
4	$\neg E$	Premise

Line	Formula	Justification
5	$\neg(B \supset D)$	Modus Tollens (3, 4)
6	$\neg B \supset \neg D$	De Morgan's Law (5)
7	$\neg B$	Elim (6)
8	$\neg A$	Modus Tollens (1, 7)
9	$\neg D$	Elim (6)
10	$\neg C$	Modus Tollens (2, 9)
11	$\neg A \supset \neg C$	Intro (8, 10)

Conclusion: We have proven that $\neg A \supset \neg C$, i.e., you did not eat carefully and you did not exercise regularly.

Slide 15: Two Special Inference Rules (continued)

- Deductive Theorem:

$$p \vdash r, p \vdash s \vdash p \implies (r \wedge s)$$
- Reductio ad absurdum:

$$p \vdash r, p \vdash \neg s \vdash p \vdash \neg r$$

Slide 16: Soundness and Completeness

- Sound: Valid argument with true premises.
- Complete: Derives any sentence entailed by premises.

Slide 17: Formal Proofs of Natural Language Arguments

- Steps:
 - Identify atomic propositions.
 - Formalize argument in logic.
 - Check for invalidity.
 - Attempt proof.

Slide 18: Example - Travel

- Argument:
 ... Therefore, if my neighbours claim to be impressed then they are just pretending.

Slide 19: Example - Travel (continued)

- Formalize argument:

$$p \implies q, \neg p \implies \neg r, \neg q \dots \neg r$$
- Proof:

$$\dots \neg p \implies r$$

Slide 20: Example - Nutrition

- Argument:

... Therefore, you did not eat carefully and you did not exercise regularly.

Slide 21: Example - Nutrition (continued)

- Formalize argument: $A \implies B, C \implies D, B \vee D \implies E, \neg E \dots \neg A \wedge \neg C$
- Proof: $\dots \neg A \wedge \neg C$

Slide 22: Application to Software Engineering

- Questions about software specifications and claims are arguments.

Slide 23: Reading and References

- Russell and Norvig, Artificial Intelligence (4th Edition)
- Nissanke, Introductory Logic and Sets for Computer Scientists
- Gray, Logic, Algebra and Databases